

Update on LISA: A Linked Slow-Start Algorithm for MPTCP

draft-barik-mptcp-lisa-01



plan: ICCRG

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Context

- This is a simple algorithm that avoids a multiple-IW burst mixed with ongoing SS
 - Reduces loss, improves Flow Completion Time (FCT)
 - In emulation and in real-life testbed (NorNet)
- Published at ICC

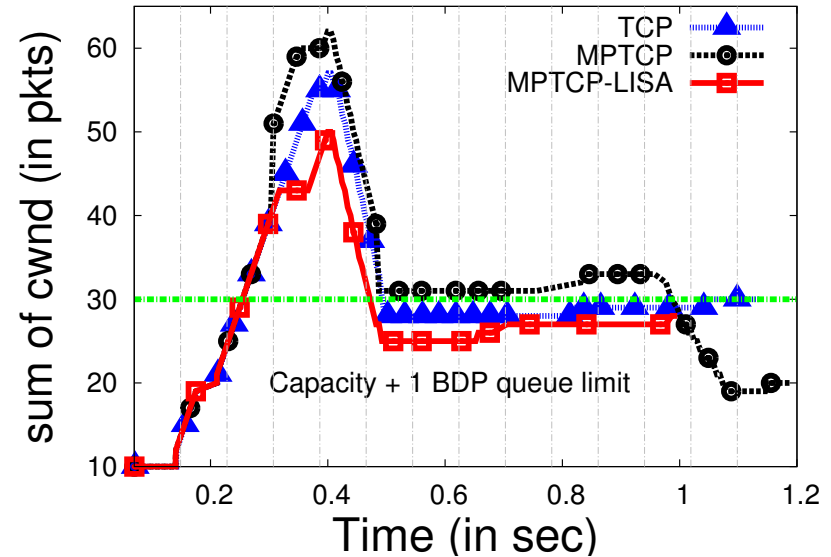
Runa Barik, Michael Welzl, Simone Ferlin, Ozgu Alay: "LISA: A Linked Slow-Start Algorithm for MPTCP", IEEE ICC 2016, Kuala Lumpur, 23-27 May 2016.
- Presented to MPTCP at IETF 94
 - Feedback: show results with different BDP
- Presented to MPTCP at IETF 96
 - Very little time, doesn't fit charter

Context /2

- Decision: discuss it in ICCRG
- If MPTCP charter updated to allow congestion control changes, bring it back there
- Else publish it via ICCRG as Experimental IRTF doc
 - Need reviews!

What's the problem?

- No coupling in SS
 - Does not seem necessary: subflow 1 doubles, subflow 2 doubles → aggregate doubles
 - However, IW of new subflows starting makes this different
 - New subflows typically start at the same time, while SS is already ongoing



2 subflows, shared bottleneck:
2.5Mbps, 70ms
transfer size: 300KB

- E.g. 8 subflows = SS + IW70

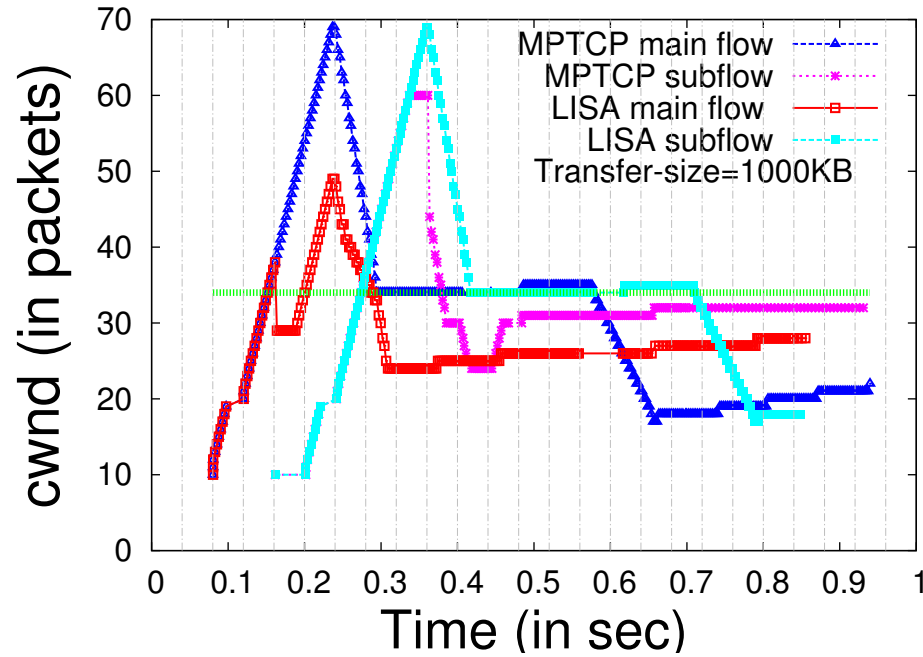
Short overview

- LISA tries to be no more aggressive than one TCP during SS
- Basic approach in LISA:
 - From subflows in SS, select subflow with maximum sending rate (`max_subflow`)
 - From `max_subflow`'s `cwnd`, take between 3 and 10 packets as “packet credit” to give `new_subflow` as IW
 - `Max_subflow` does not increase `cwnd` for $(\text{packets_inflight} - \text{cwnd})$ ACKs
 - If no `max_subflow`, IW of `new_subflow` = 10

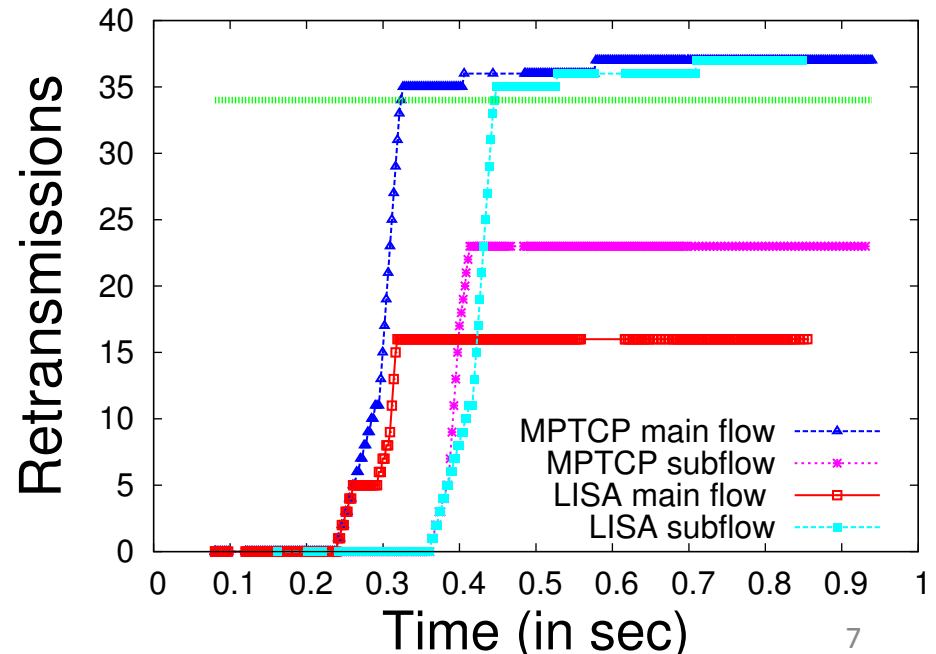
Previously shown results

- **Retransmissions & completion time**, depending on:
 - File size: problem does not occur with very small files (need to send multiple IWs), and becomes less relevant for very large files
 - Queue length: whether queue exceeded heavily (many losses & retransmissions → long FCT) or not is a matter of luck: depends on how ($\# \text{ subflows} * \text{IW}$) aligns with (BDP+queue). By reducing the burst, LISA sometimes helps
 - Number of subflows: more subflows make the problems worse and the effect of LISA better
 - RTT, Bandwidth: varying BDP has similar effect to varying queue length

Example result: non-shared BN

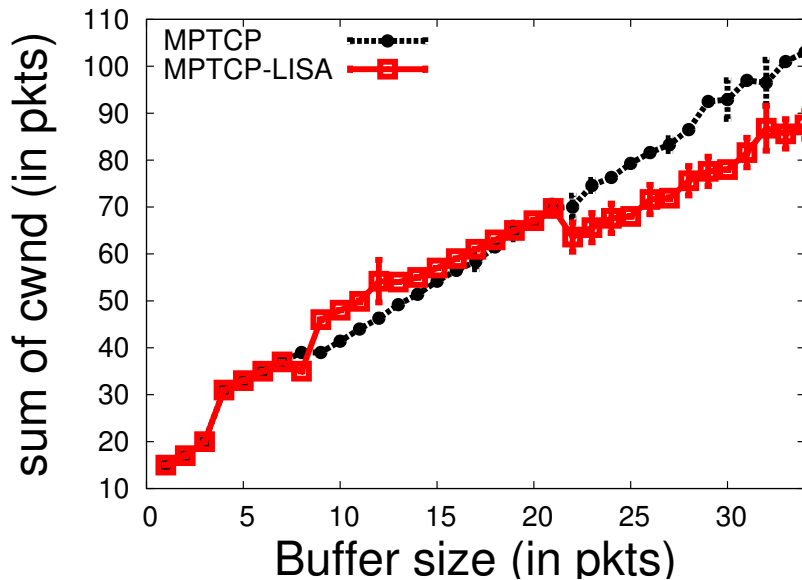
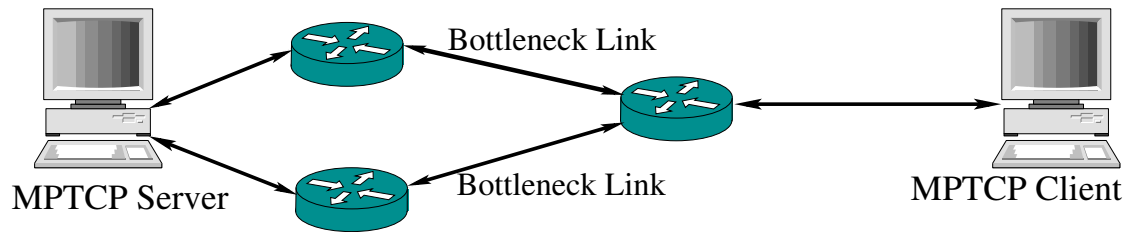


This is a “lucky-BDP” case...

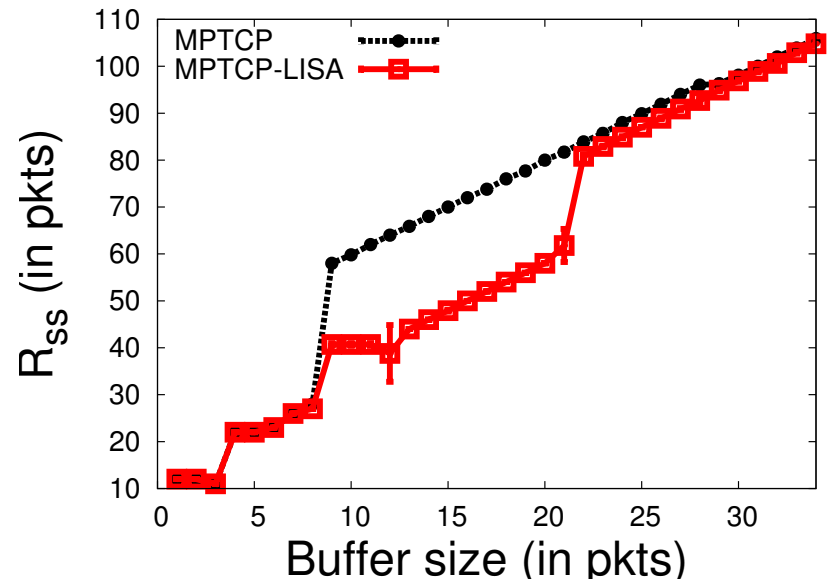


Non-shared bottleneck

5Mbps, 40ms, 2 subflows



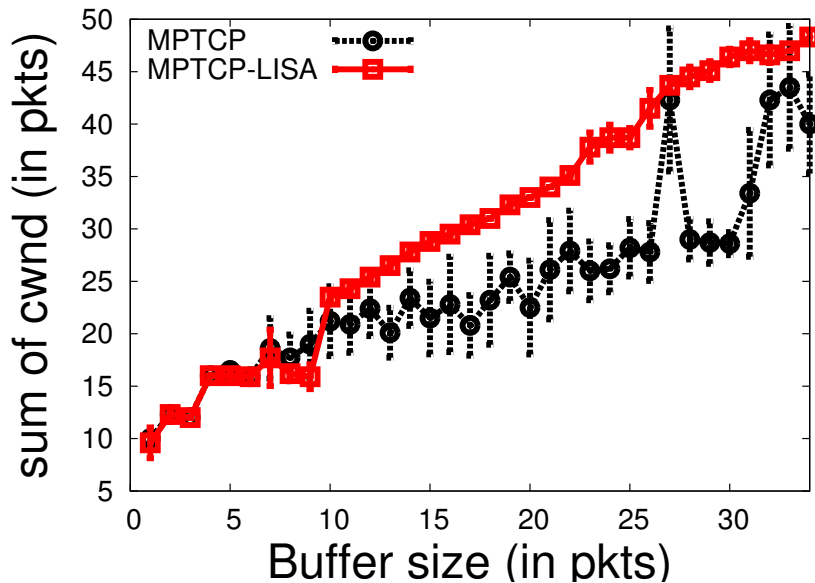
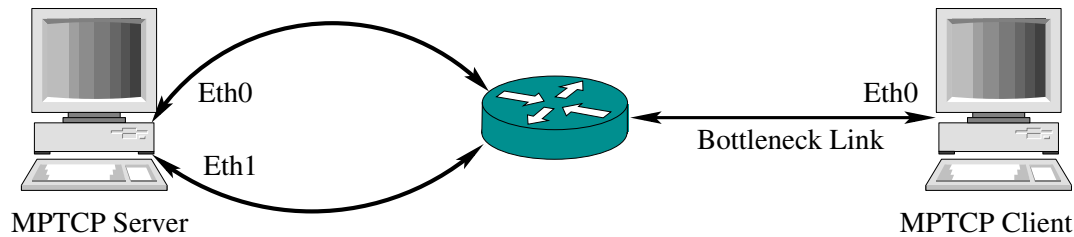
At the end of slow start (after loss recovery)



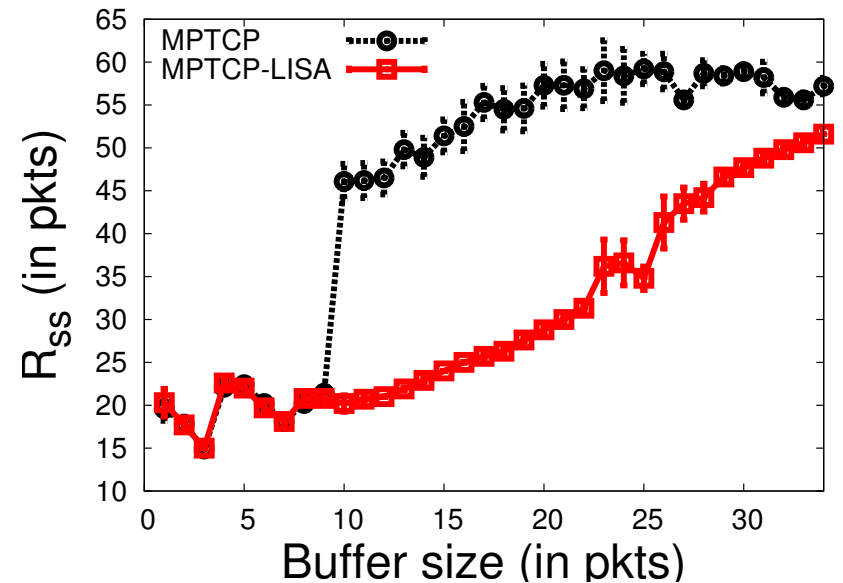
Retransmissions until then

Shared Bottleneck

5Mbps, 40ms, 2 subflows



At the end of slow start (after loss recovery)



Retransmissions until then

Next steps

- LISA paper, draft, presentations, Linux kernel patch available from:
<http://heim.ifi.uio.no/~runabk/lisa/>
- Please read & comment!